

Fast Inverse Square Roots (0x5F3759df)



Sign only defined for positive values so $S_x = 0$

$$\therefore f_x = \left(1 + \frac{M_x}{2^{n-1}}\right) 2^{E_x - (2^{b-1} - 1)}$$

For simplicity, let $L = 2^{n-1}$, which normalises the mantissa M to $0 \leq M < L$

& let $B = 2^{b-1} - 1$, which is the exponent bias.

$$f_x = \left(1 + \frac{M_x}{L}\right) 2^{E_x - B} \quad (1)$$

Looking for $y = \frac{1}{\sqrt{x}} = x^{-\frac{1}{2}}$

Let f_x & f_y be floating point representations of x & y respectively.

$$\therefore f_y = f_x^{-\frac{1}{2}} \quad \text{ignoring errors introduced by floating pt. approx.}$$

$$\log_2 f_y = \log_2 (f_x^{-\frac{1}{2}})$$

$$\log_2 f_y = -\frac{1}{2} \log_2 f_x$$

$$\log_2 \left(\left(1 + \frac{M_y}{L}\right) 2^{E_y - B} \right) = -\frac{1}{2} \log_2 \left(\left(1 + \frac{M_x}{L}\right) 2^{E_x - B} \right)$$

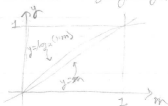
$$\log_2 \left(1 + \frac{M_y}{L}\right) + \log_2 (2^{E_y - B}) = -\frac{1}{2} \left[\log_2 \left(1 + \frac{M_x}{L}\right) + \log_2 (2^{E_x - B}) \right]$$

$$\log_2 \left(1 + \frac{M_y}{L}\right) + E_y - B = -\frac{1}{2} \log_2 \left(1 + \frac{M_x}{L}\right) - \frac{1}{2} E_x + \frac{1}{2} B$$

$$\log_a(1 + \frac{M_x}{L}) + E_y = -\frac{1}{2} \log_a(1 + \frac{M_x}{L}) - \frac{1}{2} E_x + \frac{3}{2} \mathcal{E}$$

$$\log_a(1 + \frac{M_x}{L}) + E_y = -\frac{1}{2} \log_a(1 + \frac{M_x}{L}) - \frac{1}{2} (E_x - 3\mathcal{E}) \quad (2)$$

Now consider the binary log: $\log_2(1+m)$ for $0 \leq m \leq 1$:



We see that $\log_2(1+m) \geq m$ for $0 \leq m \leq 1$.
More specifically, $\log_2(1+m) = m + \mathcal{O}_m$ for $0 \leq m \leq 1$ and since small error from \mathcal{O}_m (3)

Substituting (3) into (2), we have:

$$\frac{M_x}{L} + \mathcal{O}_y + E_y = -\frac{1}{2} \left(\frac{M_x}{L} + \mathcal{O}_x \right) - \frac{1}{2} (E_x - 3\mathcal{E})$$

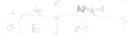
$$\frac{M_x}{L} + E_y = -\frac{1}{2} \frac{M_x}{L} - \frac{1}{2} \mathcal{O}_x - \mathcal{O}_y - \frac{1}{2} (E_x - 3\mathcal{E})$$

$$E_y L + M_y = -\frac{1}{2} M_x - L \left(\frac{1}{2} \mathcal{O}_x + \mathcal{O}_y \right) - \frac{1}{2} L (E_x - 3\mathcal{E})$$

$$= -\frac{1}{2} M_x - L \left(\frac{1}{2} \mathcal{O}_x + \mathcal{O}_y \right) - \frac{1}{2} E_x L + \frac{3}{2} \mathcal{E} L$$

$$E_y L + M_y = L \left(\frac{3}{2} \mathcal{E} - \left(\frac{1}{2} \mathcal{O}_x + \mathcal{O}_y \right) \right) - \frac{1}{2} (E_x L + M_x) \quad (4)$$

Now let's take a quick look at an n-bit floating point register again:



Cast to an integer type, then: $E 2^{N-b-1} + M = EL + M$

So if we take I_x and I_y to be the floating point values $f_x + f_y$, respectively, cast to integers, then we have

$$I_x = E_x L + M_x, \quad I_y = E_y L + M_y. \quad (5)$$

So from (5), (4) turns into

$$I_y = L \left(\frac{3}{2} B - \left(\frac{1}{2} \alpha_x + \alpha_y \right) \right) - \frac{1}{2} I_x$$

or more simply

$$I_y = R - \frac{1}{2} I_x$$

$$\text{where } R = L \left(\frac{3}{2} B - \left(\frac{1}{2} \alpha_x + \alpha_y \right) \right)$$

And that's the basic technique: take floating point x as an integer, divide it by 2 (right shift by 1), & subtract it from our magic integer R . The guess I_y , if you take that as a floating point you get y .

The error comes from the fact that α_x & α_y are for specific values of x (& therefore of y).

So we need to pick α_x & α_y as best as we can, & it will be an approximation for most values of x .

For single precision IEEE floating point, $N=32$, $e=8$

$$L = 2^{23}$$

$$B = 127$$

$$R = 2^{23} \left(\frac{3}{2} (B) - \left(\frac{1}{2} \alpha_x + \alpha_y \right) \right)$$